

The Yoneda Lemma

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1. Recap
 - A. Functors
 - B. Natural Transformations
2. Representable Functor
3. The Mini Yoneda Lemma for Type Theorists
4. The Yoneda Lemma

$$F: \mathbb{C} \longrightarrow \mathbb{D}$$

$$F: \text{Obj}(\mathbb{C}) \longrightarrow \text{Obj}(\mathbb{D}) \quad x \in \mathbb{C} \quad Fx$$

$$F: \text{Arr}(\mathbb{C}) \longrightarrow \text{Arr}(\mathbb{D}) \quad f: \mathbb{C} \quad Ff$$

$$F: \mathbb{C}(A, B) \longrightarrow \mathbb{D}(FA, FB)$$

$$\varphi: F \longrightarrow G$$

$$x \in \mathbb{C} \quad \varphi_x: Fx \longrightarrow Gx$$

$$\begin{array}{ccc}
 FX & \xrightarrow{p_x} & GX \\
 Ff \downarrow & & \downarrow Gf \\
 FY & \xrightarrow{p_y} & GY
 \end{array}$$

$$\varphi_N : \text{List } \underline{N} \longrightarrow \text{Maybe } \underline{N}$$

$$\varphi_N [] = \text{Nothing}$$

$$\varphi_N (x:xs) = \text{Just } (x+1)$$

$$\varphi_B : \text{List } B \longrightarrow \text{Maybe } B$$

$$\varphi_B [] = \text{Nothing}$$

$$\varphi_B (x:xs) = \text{Just } (\text{flip } x)$$

$$\begin{array}{ccc}
 N \xrightarrow{f} B & & \\
 \left(\begin{array}{ccc}
 \text{List } N & \xrightarrow{\varphi_N} & \text{Maybe } N \\
 \downarrow \text{List } f & & \downarrow \text{Maybe } f \\
 \text{List } B & \xrightarrow{\varphi_B} & \text{Maybe } B
 \end{array} \right) & \xrightarrow{\text{Just } (f(n))} & \text{Just } (f(n)) \\
 & & \text{Just } (\text{flip } (f(n)))
 \end{array}$$

$$f(n+1) = \text{flip}(f(n))$$

$$f: \mathbb{N} \rightarrow \mathbb{B}$$

$$f\ n = \text{True}$$

Representable

$\mathbf{HauSet}$ extends to a functor

$$\mathbb{C}(-, -) : \mathbb{C}^{\mathcal{P}} \times \mathbb{C} \rightarrow \mathbf{Set}$$

A functor $F : \mathbb{C} \rightarrow \mathbf{Set}$ is representable whenever $\exists A \in \mathbf{Obj}(\mathbb{C})$ s.t.

$$F \cong \mathbb{C}(A, -)$$

$$\mathbf{List} \cong \mathbf{Set}(\mathbb{N}, -)$$

is it true that for all A $\mathbf{List} A \cong \mathbf{Set}(\mathbb{N}, A)$

$$\mathbf{List} A \not\cong \mathbb{N} \rightarrow A$$

In Set

$$\text{Str } A \cong \mathbb{N} \rightarrow A$$

$$xs \mapsto \lambda n. \text{head (drop } n \text{ xs)}$$

$$f.\text{succ} \leftarrow \lambda f$$

(Mini) Yoneda Lemma for Type Theorists

$$\frac{\Gamma \vdash t : A}{\Gamma \vdash R(t) : B}$$

Task is find an interpretation for R
given a morphism $A \xrightarrow{f} B$

$$\in \mathcal{C}(\Gamma, B)$$

$$[\Gamma \vdash R(t) : B] : [\Gamma] \rightarrow [B]$$

$$\boxed{\begin{array}{c} \text{HptInd} \\ [\Gamma] \rightarrow [A] \end{array}}$$

$$\frac{\Gamma \xrightarrow{t} A}{\Gamma \xrightarrow{R(t)} B}$$

$$A \xrightarrow{f} B$$

Mini Yoneda $\mathcal{C}$ locally small

$$\mathcal{C}(A, B) \cong \mathcal{C}(-, A) \xrightarrow{\circ} \mathcal{C}(-, B)$$

PO

$$\cong \boxed{\forall \Gamma. \mathbb{C}(\Gamma, A) \rightarrow \mathbb{C}(\Gamma, B)}$$

$$\frac{[\Gamma \vdash t : A] \in \mathbb{C}(\Gamma, A)}{[\Gamma \vdash R(t) : B] \in \mathbb{C}(\Gamma, B)}$$

$$\mathbb{C}(A, B) \xrightarrow[\psi]{\varphi} \mathbb{C}(-, A) \rightarrow \mathbb{C}(-, B) \quad \neq$$

$$\psi(f)(g : \Gamma \rightarrow A) = \underbrace{\lambda \alpha. f(g(\alpha))}_{\Gamma \rightarrow B} = f \cdot g$$

$$\psi(\alpha) = \alpha_A (id_A) \rightarrow \mathbb{C}(A, A) \rightarrow \mathbb{C}(A, B)$$

$$\frac{\Gamma \vdash t : \boxed{A}}{\Gamma \vdash R(t) : \boxed{B}} \quad \neq$$

Yoneda $F: \mathcal{C}^{\text{op}} \rightarrow \text{Set}$

$$FC \cong \mathcal{C}(-, C) \xrightarrow{\cdot} F$$

for example $F = \mathcal{C}(-, B)$

$$\underline{FA} \cong \mathcal{C}(-, A) \xrightarrow{\cdot} \mathcal{C}(-, B)$$

$$\mathcal{C}(A, B)$$